

## Formale Übersetzungsmodelle

### Task 21 (BOT<sup>2</sup> and TOP<sup>2</sup>)

Let  $\Sigma = \{\sigma^{(2)}, \gamma^{(1)}, \alpha^{(0)}, \beta^{(0)}\}$  be a ranked alphabet. Consider the bu-tt  $B = (Q_B, \Sigma, \Sigma, F, R_B)$  and the td-tt  $T = (Q_T, \Sigma, \Sigma, I, R_T)$  where  $Q_B = \{*, q, q_f\}$ ,  $F = \{q_f\}$ ,  $Q_T = \{*, q\}$ ,  $I = \{*\}$ , and

$$R_B = \{ \begin{array}{l} \sigma(*(\bar{x}_1), *(\bar{x}_2)) \rightarrow *(\sigma(\bar{x}_1, \bar{x}_2)), \\ \sigma(*(\bar{x}_1), q(\bar{x}_2)) \rightarrow q_f(\bar{x}_1), \\ \gamma(*(\bar{x}_1)) \rightarrow *(\gamma(\bar{x}_1)), \\ \gamma(q(\bar{x}_1)) \rightarrow q(\gamma(\bar{x}_1)), \\ \gamma(q_f(\bar{x}_1)) \rightarrow q_f(\gamma(\bar{x}_1)), \\ \alpha \rightarrow *(\alpha), \quad \alpha \rightarrow q(\alpha), \quad \beta \rightarrow q(\beta) \end{array} \} \quad R_T = \{ \begin{array}{l} q(\sigma(\bar{x}_1, \bar{x}_2)) \rightarrow \sigma(q(\bar{x}_1), q(\bar{x}_2)), \\ *(\sigma(\bar{x}_1, \bar{x}_2)) \rightarrow \sigma(q(\bar{x}_1), *(\bar{x}_1)), \\ *(\sigma(\bar{x}_1, \bar{x}_2)) \rightarrow \sigma(*(\bar{x}_1), q(\bar{x}_1)), \\ *(\gamma(\bar{x}_1)) \rightarrow \gamma(*(\bar{x}_1)), \\ q(\gamma(\bar{x}_1)) \rightarrow \gamma(q(\bar{x}_1)), \\ *(\alpha) \rightarrow \alpha, \quad q(\alpha) \rightarrow \alpha, \quad *(\beta) \rightarrow \beta \end{array} \}$$

- (a) Identify the bottom-up and top-down specific properties of  $\tau(B)$  and  $\tau(T)$ .
- (b) Give td-tt  $T_1$  and  $T_2$  and bu-tt  $B_1$  and  $B_2$  such that  $\tau(B) = \tau(T_1); \tau(T_2)$  and  $\tau(T) = \tau(B_1); \tau(B_2)$ .

### Task 22 (Baker's theorem for BOT)

**Theorem** [Bak79, Thm. 6]. Let  $B_1$  and  $B_2$  be bu-tt. Then  $\tau(B_1); \tau(B_2) \in \text{BOT}$  if the following two conditions hold:

1.  $B_1$  is linear or  $B_2$  is deterministic;
2.  $B_1$  is nondeleting or  $B_2$  is total.

- (a) Give two bu-tt  $B'_1$  and  $B'_2$  that fulfill Condition 1 but not Condition 2 such that  $\tau(B'_1); \tau(B'_2) \notin \text{BOT}$ . Give two bu-tt  $B''_1$  and  $B''_2$  that do not fulfill Condition 1 but fulfill Condition 2 such that  $\tau(B''_1); \tau(B''_2) \notin \text{BOT}$ . For each bu-tt, use the minimum number of rules necessary.
- (b) Construct the instance  $B'$  and  $B''$  (for  $B'_1$  and  $B'_2$ , and  $B''_1$  and  $B''_2$ , respectively) of the bu-tt  $B$  defined in the proof (from the lecture) of the above theorem.
- (c) Give trees  $s', t', s'', t''$  such that
  - (i)  $\neg((s', t') \in \tau(B'_1) \circ \tau(B'_2)) \iff (s', t') \in \tau(B')$  and
  - (ii)  $\neg((s'', t'') \in \tau(B''_1) \circ \tau(B''_2)) \iff (s'', t'') \in \tau(B'')$ .
- (d) Prove the following corollary:
 

**Corollary.** Let  $B_1$  and  $B_2$  be bu-tt. Then  $\tau(B_1); \tau(B_2) \in \text{BOT}$  if  $B_1$  is linear or  $B_2$  is deterministic.
- (e) Apply the above corollary to  $B'_1$  and  $B'_2$  from Task 22 (a).

## **References**

- [Bak79] Brenda S. Baker. Composition of top-down and bottom-up tree transductions. *Information and Control*, 41(2):186–213, 1979.