

Formale Baumsprachen

Task 4 (bu-det fta)

Let $\Sigma = \{\sigma^{(2)}, \alpha^{(0)}, \beta^{(0)}\}$ and $\Delta = \{\sigma^{(2)}, \gamma^{(1)}, \alpha^{(0)}\}$ be ranked alphabets. Give deterministic bu-ta \mathcal{A}_1 , \mathcal{A}_2 , and \mathcal{A}_3 that recognize L_1 , L_2 , and L_3 , respectively, where

- (a) $L_1 = \{\xi \in T_\Sigma \mid \xi \text{ contains at least one } \alpha \text{ and one } \beta\}$,
- (b) $L_2 = \{\xi \in T_\Sigma \mid \xi \text{ contains an even number of } \alpha \text{ symbols}\}$, and
- (c) $L_3 = \{\sigma(t_1, \sigma(t_2, \dots, \sigma(t_n, \alpha) \dots)) \in T_\Delta \mid n \in \mathbb{N}, t_1, \dots, t_n \in T_{\{\gamma^{(1)}, \alpha^{(0)}\}}\}$.

Task 5 (string automata)

Recall the concept of string automata. Let Σ be an alphabet and $\# \notin \Sigma$. We define the ranked alphabet $\Sigma_\# = \Sigma_\#^{(0)} \cup \Sigma_\#^{(1)}$ where $\Sigma_\#^{(0)} = \{\#\}$ and $\Sigma_\#^{(1)} = \Sigma$. Moreover, we define the $\Sigma_\#$ -algebra (Σ^*, θ) where $\theta(\#) = \varepsilon$ and $\theta(a)(w) = wa$ for every $a \in \Sigma$ and $w \in \Sigma^*$.

- (a) Show that Σ^* is initial in the class of $\Sigma_\#$ -algebras.
- (b) We consider $\Sigma = \{a, b\}$ and the language $L = \{a^n b^m \mid n, m \in \mathbb{N}\}$. Sketch the diagram of a total deterministic finite-state automaton accepting L and model the transition table using a finite $\Sigma_\#$ -algebra Q . How can we interpret the uniquely determined homomorphism $h: \Sigma^* \rightarrow Q$?
- (c) Convince yourself that any total deterministic finite-state automaton can be modeled as a quadruple $\mathcal{A} = (Q, \Sigma, \theta, F)$ where (Q, θ) is a finite $\Sigma_\#$ -algebra and $F \subseteq Q$. Define the language accepted by \mathcal{A} using the homomorphism $h: \Sigma^* \rightarrow Q$.

Task 6 (regular tree grammars)

Let $\Sigma = \{\sigma^{(2)}, \gamma^{(1)}, \alpha^{(0)}\}$ be a ranked alphabet. Give regular tree grammars G_1 and G_2 with

- (a) $L(G_1) = \{\xi \in T_\Sigma \mid \xi \text{ contains exactly one } \sigma\}$ and
- (b) $L(G_2) = \{\xi \in T_\Sigma \mid \xi \text{ contains the pattern } \sigma(_, \gamma(_)) \text{ at least twice}\}$.